



ALMA MATER STUDIORUM
UNIVERSITÀ DI BOLOGNA



Babele è stata davvero una maledizione? La molteplicità dei linguaggi di programmazione

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Matematica, Fisica e Informatica nel secolo XX:
L'ossessione della totalità
Accademia delle Scienze di Torino, 19 marzo 2024

What is a programming language?

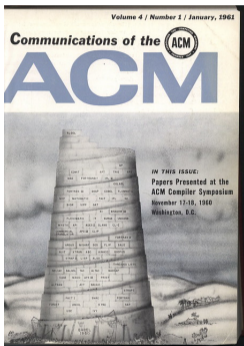
An artificial language used to write instructions that can be translated into machine language and then executed by a computer.

[THE AMERICAN HERITAGE SCIENCE DICTIONARY, ©2011]

How many programming languages are there?

Thousands... Several dozen currently in use

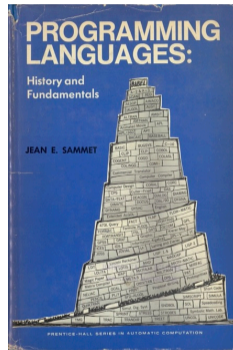
The Babel of programming languages



1961



1964



1969

Babel vs Universal: Saul Gorn (1912-1992)

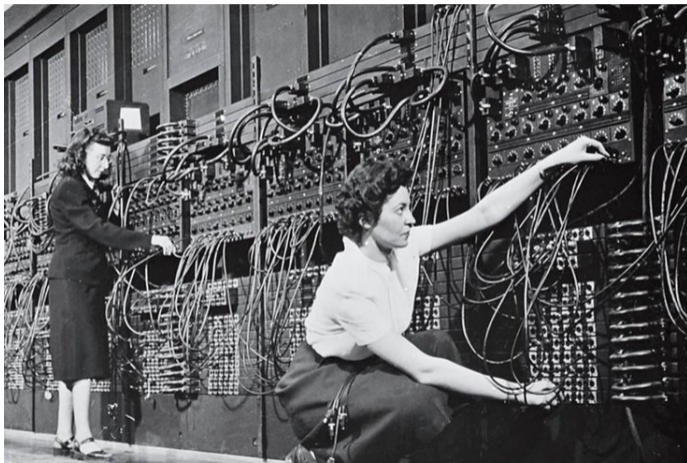


A simple enough “universal code” to be used by “computers, data processors, production engineers, traffic controllers, or administrators of large companies.”

[Planning universal semi-automatic coding, 1954]

The first proliferation: machines

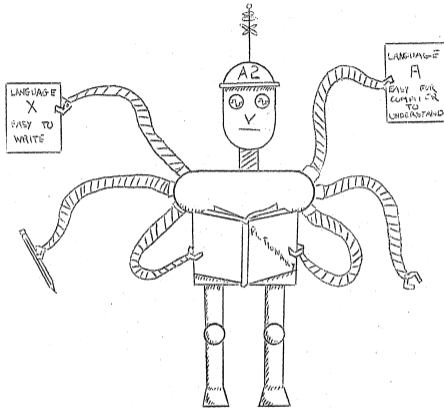
ENIAC: 1947



One machine, one code

Technology, bound to the **specific** machine

Translation: Grace M. Hopper



TRANSLATION PHASE

H/S 10/2/53

[Digital computer advanced coding techniques, ©MIT, 1954]

Towards the language metaphor

When Technology Became Language

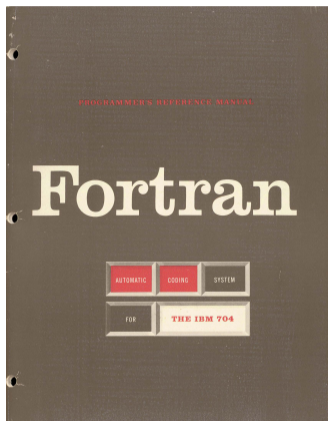
The Origins of the Linguistic Conception of Computer Programming, 1950–1960

**DAVID NOFRE, MARK PRIESTLEY, and
GERARD ALBERTS**

Technology and Culture, Vol 55, January 2014.

The second proliferation: domains

Algebraic expressions: FORTRAN



©IBM, 1956

Algebraic expressions, e.g.

$$D=(A+B)*C-\sin(A*C+2)$$

translated into **efficient** object programs

Simple control structures: FORTRAN

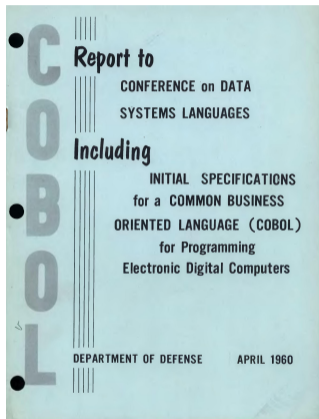
A DO Nest with Exit and Return

Given an $N \times N$ square matrix A , to find those off-diagonal elements which are symmetric and to write them on binary tape.

FOR COMMENT		CONTINUATION	FORTRAN STATEMENT	IDENTIFICATION				
STATEMENT NUMBER	1			5	6	7	71	73
			REWIND 3					
			DO 3 I = 1,N					
			DO 3 J = 1,N					
			IF(A(I,J)-A(J,I)) 3,20,3					
	3		CONTINUE					
			END FILE 3					
			MORE PROGRAM					
	20		IF(I-J) 21,3,21					
	21		WRITE TAPE 3,I,J, A(I,J)					
			GO TO 3					

©IBM, 1956

Records and English: COBOL



©US Dept of Defense, 1960

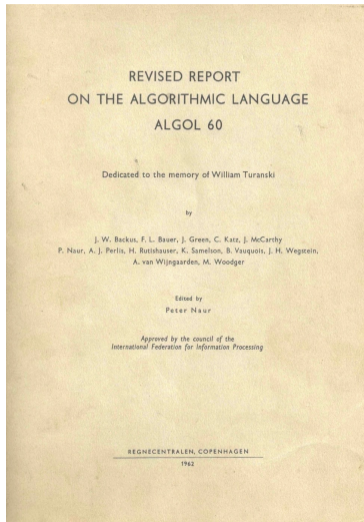
Collections of non-numerical data
English words

Records and English: COBOL

```
FD  CUSTOMER-FILE
    RECORD CONTAINS 45 CHARACTERS.
01  CUSTOMER-RECORD.
    05  CUSTOMER-NAME.
        10  LAST-NAME          PIC X(17).
        10  FILLER             PIC X.
        10  INITIALS          PIC XX.
    05  PART-ORDER.
        10  PART-NAME          PIC X(15).
        10  PART-COLOR        PIC X(10).
```

```
IF LAST-NAME = PART-NAME GO TO PARAGRAPH 1 ELSE
MOVE PART-NAME TO LAST-NAME
```

Algorithms: ALGOL



Universal language for algorithm exchange

International committee

Algorithms: ALGOL

ALGORITHM 64

QUICKSORT

C. A. R. HOARE

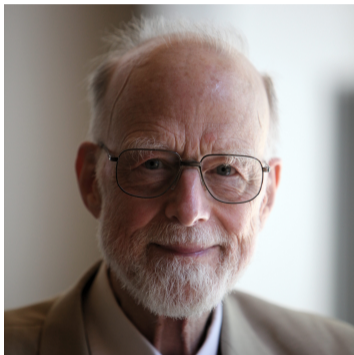
Elliott Brothers Ltd., Borehamwood, Hertfordshire, Eng.

```
procedure quicksort (A,M,N); value M,N;  
      array A; integer M,N;
```

comment Quicksort is a very fast and convenient method of sorting an array in the random-access store of a computer. The entire contents of the store may be sorted, since no extra space is required. The average number of comparisons made is $2(M-N) \ln(N-M)$, and the average number of exchanges is one sixth this amount. Suitable refinements of this method will be desirable for its implementation on any actual computer;

```
begin      integer I,J;  
          if M < N then begin partition (A,M,N,I,J);  
                                quicksort (A,M,J);  
                                quicksort (A, I, N)  
          end  
end      quicksort
```

Rich data types: Tony Hoare



C.A.R. Hoare, 1934-

Modelling tool:

In the simulation of complex situations in the real world, it is necessary to construct in the computer analogues of the objects of the real world

[Hoare, Record handling, 1965]

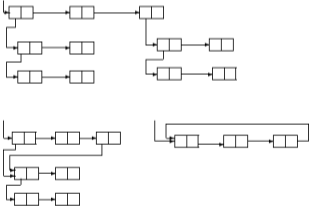
Symbolic structures: LISP



John McCarthy, 1927-2011

Recursive Functions of Symbolic Expressions and Their Computation by Machine, Part I

Communications of the ACM Volume 3, April 1960



Symbolic structures: LISP

```
eval[e; a] = [  
  atom[e] → cdr[assoc[e;a]];  
  atom[car[e]] → [eq[car[e]; QUOTE] → cadr[e]];  
                  eq[car[e]; COND] → evcon[cdr[e];a];  
                  T → apply[car[e]; evlis[cdr[e];a];a];  
  T → apply[car[e]; evlis[cdr[e];a];a]
```

Different domains, different machines

Scientific: FORTRAN, on IBM the 7090 and the IBM 1620

Business: COBOL, on the IBM 7080 and the IBM 1401

Real-time: JOVIAL, on the IBM 7750 and IBM 7950 (Harvest)

Abstraction over the machine: Hoare

It was a firm principle of our implementation that the results of any program, even erroneous, should be comprehensible without knowing anything about the machine or its storage layout.

[Hoare, personal communication, 2014]

Taking stock, 1

1. PLs do not give instructions to the (physical) machine: they **hide** it.
2. PLs are sets of **abstraction mechanisms**,
 - over control (structured control, procedures)
 - and data (data types).
3. Programs are abstract, computational **models** of “the real world” (cf Hoare).

The myth of the total language

IBM System/360, 1964 ff



The ultimate language

One machine, one language: for all

- ▶ users need to learn only one language
- ▶ only one compiler to be maintained
- ▶ programs could be easily shared

“A universal programming language that would meld and displace FORTRAN and COBOL”

[Brooks and Shustek, 2015]

PL/I: Some design choices

- ▶ Anything goes
 - “If a particular combination of symbols has a reasonably sensible meaning, that meaning will be made official”*
- ▶ Full access to machine and operating system
- ▶ Cater to the novice

[G. Radin, H.P. Rogoway. NPL: Highlights of A New Programming Language. CACM 8(1), 1965]

PL/I: An inconsistent model

PL/I: An inconsistent model



PL/I: defeat

Different domains raise different classes of problems that require different sets of representations.

PL/I was **designed** in order to forget about such peculiarities.

Other driving forces: correctness

Algol's research programme:

a (Kuhn) paradigm for programming language design, and
correct software development.

A language for the new science

Ada: 1977 ff

Designed *for* the US Department of Defence:

- ▶ concurrency
- ▶ real-time
- ▶ embedded computing
- ▶ life-critical applications
- ▶ reliability
- ▶ formal definition
- ▶ simplicity

Ada, the last total language

Jean Ichbiah (Ada's main designer):

In ten years from now [scil. 1979-80], only two programming languages will remain: Ada and Lisp.

[according to Rosen, The Ada paradox(es), Ada Letters 24, 2009]

Another attempt to universality

Total language



Total language

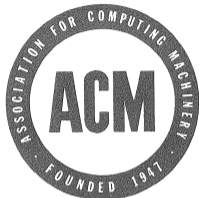


vs.

Extensible language



Extensible languages



An informal monthly publication of the *Special Interest Group on Programming Languages* (SIGPLAN) of the *Association for Computing Machinery* (ACM), incorporating the *PL/I Bulletin*, the *Snobol Bulletin*, the *Algol Bulletin*, the *LISP Bulletin*, and the *Fortran Information Bulletin* as occasional supplements.

Current SIGPLAN officers are: the Chairman, Prof. Peter Wegner, Division of Applied Mathematics, Brown University, Providence, Rhode Island 02912, telephone 401/863 2115; the Vice-Chairman, Dr. Thomas A. Standish, Aiken Computation Laboratory, Harvard University, 33 Oxford Street, Cambridge,

SIGPLAN Notices

Vol. 4, No. 8, 1969 August

SPECIAL INTEREST GROUP ON PROGRAMMING LANGUAGES

PROCEEDINGS OF THE

Extensible Languages Symposium

edited by

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sponsored by SIGPLAN

Boston, Massachusetts, 1969 May 13

Symposium Chairman: Carlos Christensen, ADR/Computer Associates

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Peter Wegner, Cornell University

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An assessment on extensible languages, 1975

Extending a simple base results often in long, thin extension cascades that are often ugly and inefficient.

[Standish, Extensibility in programming language design. AFIPS 1975]

An assessment on extensible languages, 1975

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[Standish, Extensibility in programming language design. AFIPS 1975]

Buy instead of build

Nice try, though

Total languages are *closed* (technical) objects

Extensible languages are (more) *open* (technical) objects

Taking stock, again

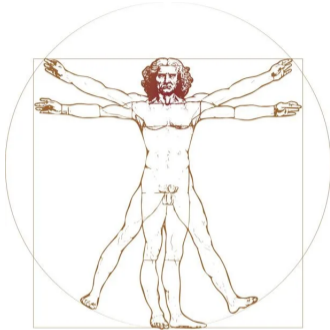
Taking stock, 1 (again)

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Taking stock, 2

4. Each PL has its own “embedded” model
5. Different domains need different models
6. External drivers

An interesting pluralism: Languages as mediators



Complete relativism?

No criteria for discernment?

Let's take the “language metaphor” seriously:
a PL is a medium for dialogue with the machine

Let's take the "language metaphor" seriously:
a PL is a medium for dialogue with the machine

The machine is a source of alienation

The criteria: reduce alienation

Gilbert Simondon



Gilbert Simondon, 1924-1989

*Les objets techniques qui
produisent le plus d'aliénation
sont ceux qui sont destinés à des
utilisateurs ignorants.*

[Du mode d'existence des objets techniques, 1958]

Open and closed technical objects

Closed technical object

- ▶ its user does not understand how and why it works
- ▶ it cannot be repaired
- ▶ it is unmodifiable
- ▶ it evokes the sacred, the untouchable

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Open technical object

- ▶ its user knows how it works, and how it could be repaired
- ▶ “to be” instead of “to appearing” (*être et ne pas paraître*)
- ▶ it shows the traces of its own evolution

“Open” programming languages

Let everyone be allowed to use a language that suits them

A language that reveals and mediates the machine within the limits, aspirations, and competences of **that** user

Such a language can reduce their alienation

Languages as extensible systems

Logo: S. Papert et al., 1967. BBN, MIT

Smalltalk: A. Kay et al., 1975. Xerox PARC

Per conoscere il mondo bisogna costruirlo

Cesare Pavese, Il mestiere di vivere. 1952

Babel?



©E. De Guzman, 2014

ἤκουον εἷς ἕκαστος τῆ ἰδίᾳ διαλέκτῳ λαλούντων αὐτῶν·

[At. 2, 6]

*Babel was the contrary of a curse.
The gift of tongues is precisely that;
a gift and benediction beyond reckoning.*

[G. Steiner. Errata. 1998 (p. 99)]

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How monotone must love-making have been in Paradise.

[(p. 102)]

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*programming**

[(p. 102)]